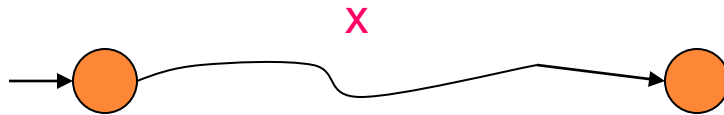


CONVERSION FROM NFA TO DFA



DFA

- For every string x , there is a **unique** path from initial state and associated with x .



x is accepted if and only if this path ends at a final state.



NFA

- For any string x , there may exist **none** or **more than one** path from initial state and associated with x .



NFA \rightarrow DFA

○ Consider an NFA $M=(Q, \Sigma, \delta, s, F)$.

○ For x in Σ^* , define

$$[x] = \{q \text{ in } Q \mid \text{there exists a path } s \rightarrow \dots \rightarrow q\}$$

○ Define DFA $M'=(Q', \Sigma, \delta', s', F')$:

$$Q' = \{ [x] \mid x \text{ in } \Sigma^* \},$$

$$\delta' ([x], a) = [xa] \text{ for } x \text{ in } \Sigma^* \text{ and } a \text{ in } \Sigma,$$

$$s' = [s],$$

$$F' = \{ [x] \mid x \text{ in } L(M) \}$$



CONSTRUCTION OF M'

Special Case: M has no ϵ -move.

- $[\epsilon] = \{s\}$
- Suppose $[x]$ is known. How to get $[xa]$ for a in Σ ?



FROM [X] TO [XA]

○ $[xa] = \{ p \mid \text{there exists a path } q \rightarrow p \text{ for some } q \text{ in } [x] \}$

$= \{ p \mid \text{there exists } q \text{ in } [x],$
 $q \rightarrow p \}$

$= \bigcup_{q \in [x]} \delta(q, a)$

$\xrightarrow{\text{a}}$
path

$\xrightarrow{\text{a}}$
edge



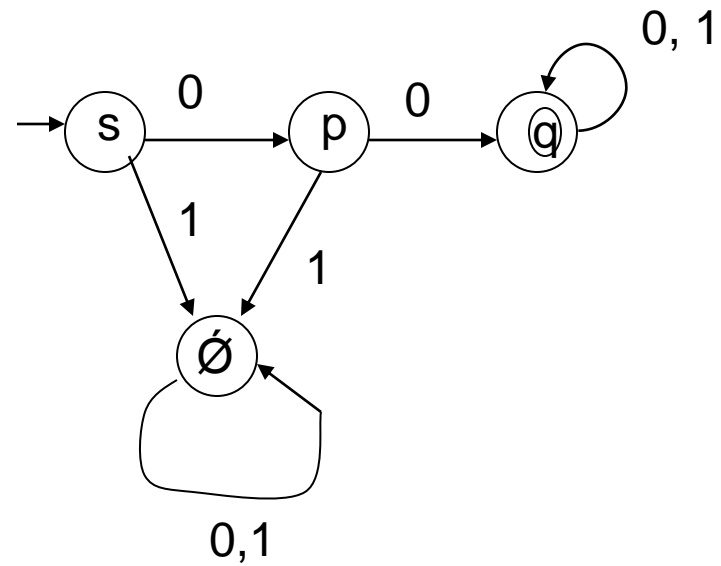
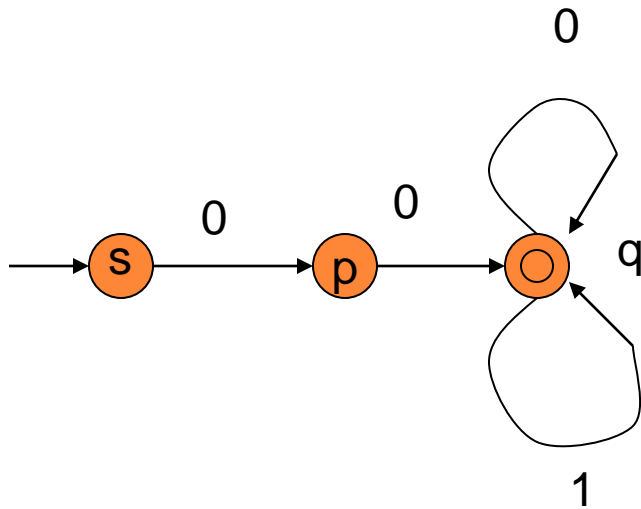
CONSTRUCTION OF M'

- $F' = \{[x] \mid x \text{ in } L(M)\}$
 $= \{[x] \mid [x] \cap F \neq \emptyset\}$



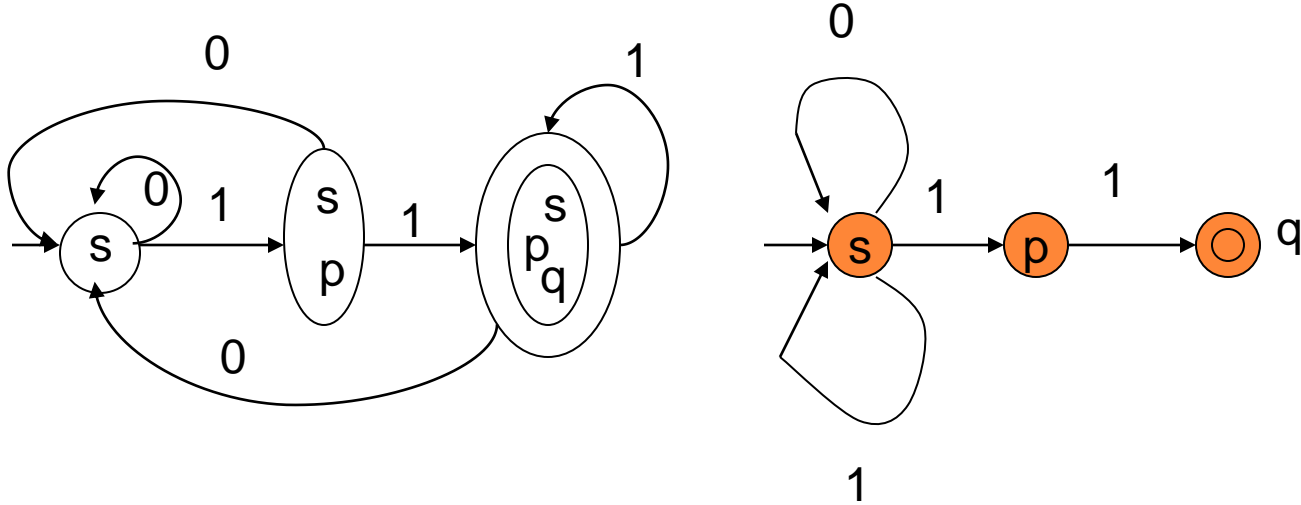
EXAMPLE 1

- Construct DFA to accept $00(0+1)^*$



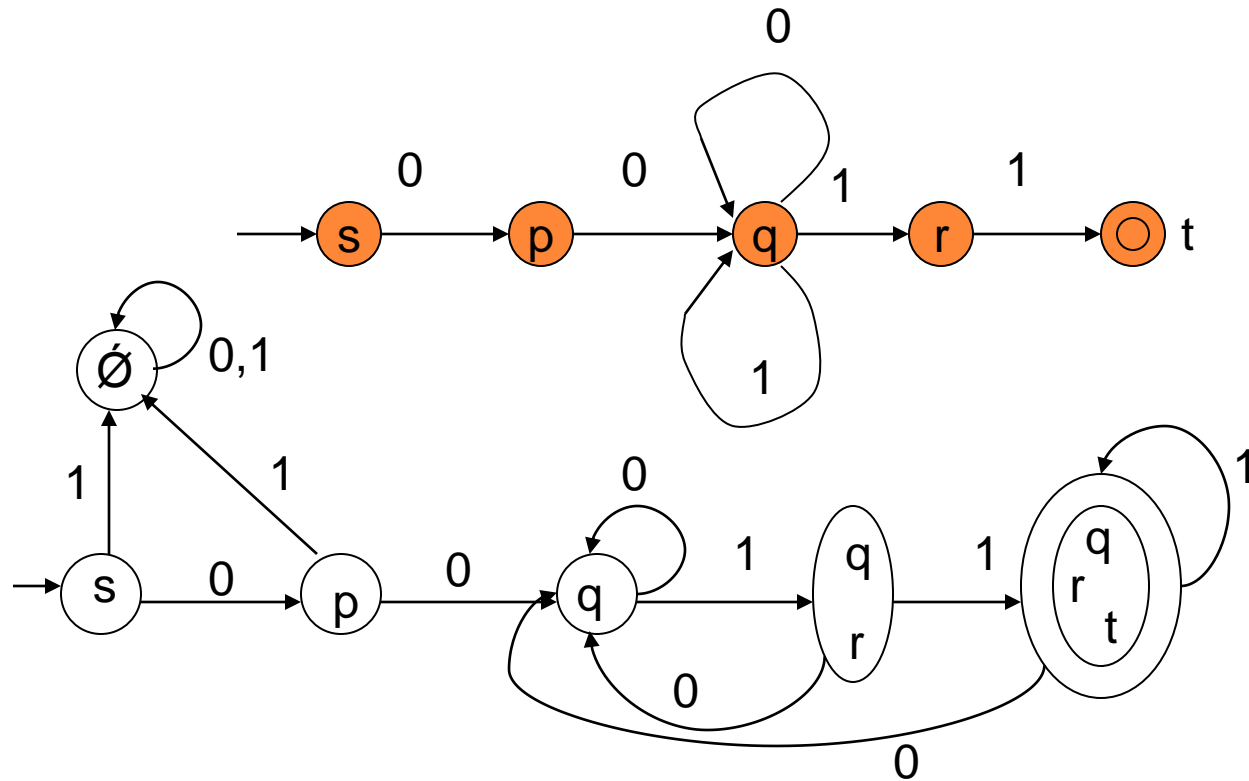
EXAMPLE 2

- Design DFA to accept $(0+1)^*11$



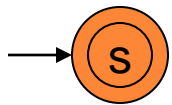
EXAMPLE 3

- Design DFA to accept $00(0+1)^*11$

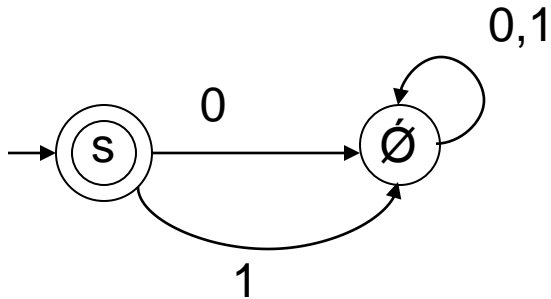


EXAMPLE 4

- Construct DFA M for $L(M)=\epsilon$.



Is this a DFA?

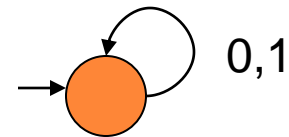
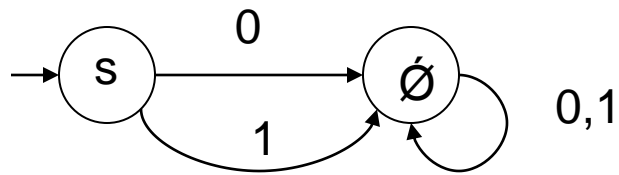


EXAMPLE 5

- Construct DFA M for $L(M)=\emptyset$.



Is it a DFA?



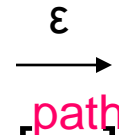
CONSTRUCTION OF M'

- For q in Q , define

ε -closure(q) = { p | there exists a path $q \rightarrow p$ }

- $[\varepsilon] = \{q \mid \text{there is a path } s \rightarrow q\}$
= ε -closure(s)

- Suppose $[x]$ is known. How to get $[xa]$ for a in Σ ?



FROM [X] TO [XA]

- $[xa] = \{ p \mid \text{there exists a path } q \rightarrow p \text{ for } \xrightarrow[\text{path}]{a} \text{ some } q \text{ in } [x] \}$
- $= \{ p \mid \text{there exists } q \text{ in } [x],$

$$q \rightarrow r \rightarrow p \}$$
- $= \{ p \mid \text{for some } q \text{ in } [x] \text{ and } r \text{ in } \delta(q,a),$

$$p \text{ in } \epsilon\text{-closure}(r) \}$$
- $= \bigcup_{q \text{ in } [x]} \bigcup_{r \text{ in } \delta(q,a)} \epsilon\text{-closure}(r)$

$$q \text{ in } [x] \quad r \text{ in } \delta(q,a)$$



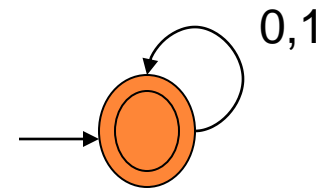
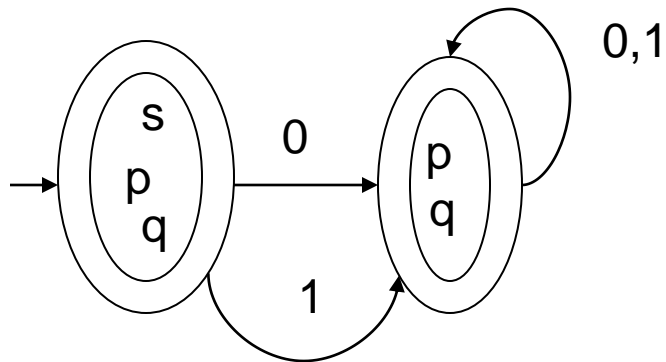
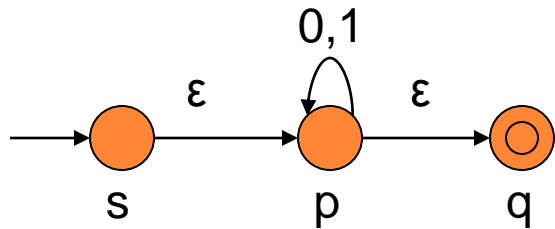
CONSTRUCTION OF M'

- $F' = \{[x] \mid x \text{ in } L(M)\}$
 $= \{[x] \mid [x] \cap F \neq \emptyset\}$



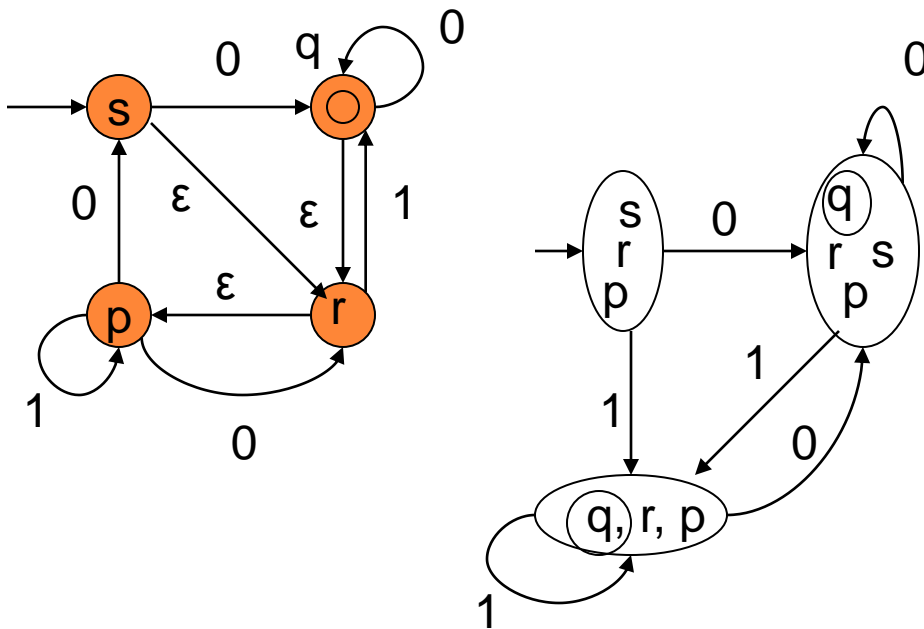
EXAMPLE 6

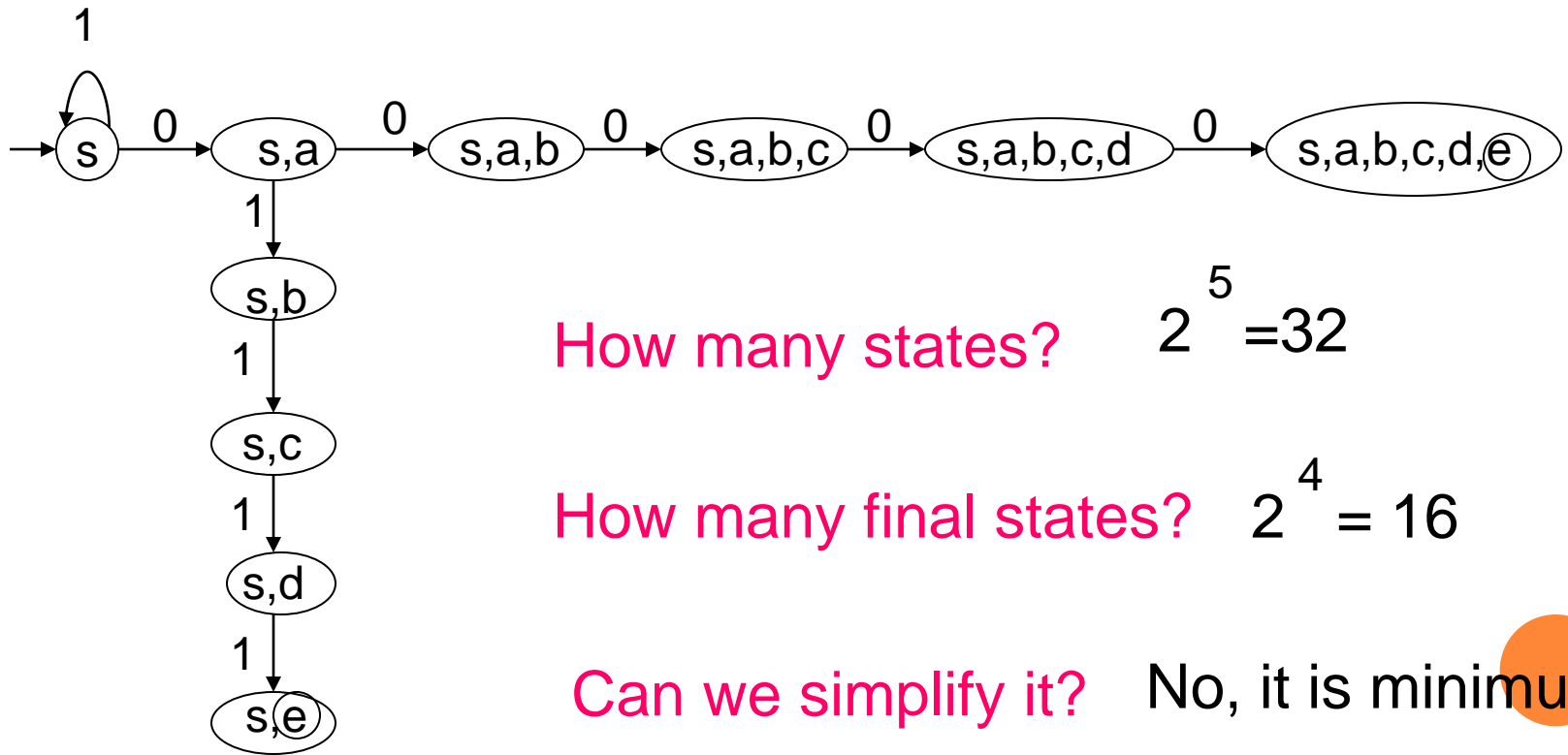
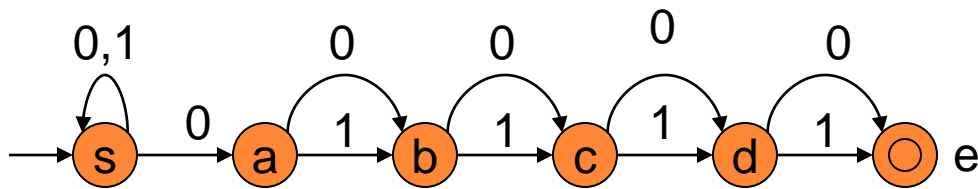
- Construct DFA M for $L(M)=(0+1)^*$.



EXAMPLE 7

- Convert the following NFA to DFA.





How many states? $2^5 = 32$

How many final states? $2^4 = 16$

Can we simplify it? No, it is minimum!